

A Dirac-type theorem for Berge cycles in random hypergraphs

Dennis Clemens

Institut für Mathematik
Technische Universität Hamburg-Harburg
Hamburg, Germany
dennis.clemens@tuhh.de

Julia Ehrenmüller

Ehrenmüller GmbH
Kempten, Germany
julia@ehrenmueller.ai

Yury Person*

Institut für Mathematik
Technische Universität Ilmenau
Ilmenau, Germany
yury.person@tu-ilmenau.de

Submitted: Mar 25, 2019; Accepted: May 13, 2020; Published: Aug 21, 2020

© The authors. Released under the CC BY-ND license (International 4.0).

Abstract

A Hamilton Berge cycle of a hypergraph on n vertices is an alternating sequence $(v_1, e_1, v_2, \dots, v_n, e_n)$ of distinct vertices v_1, \dots, v_n and distinct hyperedges e_1, \dots, e_n such that $\{v_1, v_n\} \subseteq e_n$ and $\{v_i, v_{i+1}\} \subseteq e_i$ for every $i \in [n-1]$. We prove the following Dirac-type theorem about Berge cycles in the binomial random r -uniform hypergraph $H^{(r)}(n, p)$: for every integer $r \geq 3$, every real $\gamma > 0$ and $p \geq \frac{\ln^{17r} n}{n^{r-1}}$ asymptotically almost surely, every spanning subgraph $H \subseteq H^{(r)}(n, p)$ with minimum vertex degree $\delta_1(H) \geq \left(\frac{1}{2^{r-1}} + \gamma\right) p \binom{n}{r-1}$ contains a Hamilton Berge cycle. The minimum degree condition is asymptotically tight and the bound on p is optimal up to some polylogarithmic factor.

Mathematics Subject Classifications: 05C80, 05D40

1 Introduction

Many classical theorems of extremal graph theory give sufficient optimal minimum degree conditions for graphs to contain copies of large or even spanning structures. Lately it

*YP was supported by DFG grant PE 2299/1-1. An extended abstract of this paper appeared in the Proceedings of Discrete Mathematics Days 2016 (Barcelona) [7].

became popular to phrase such extremal results in terms of local resilience, where the *local resilience* of a graph G with respect to a given monotone increasing graph property \mathcal{P} is defined as the minimum number $\rho \in \mathbb{R}$ such that one can obtain a graph without property \mathcal{P} by deleting at most $\rho \cdot \deg(v)$ edges from every vertex $v \in V(G)$. For instance, using this terminology, Dirac's theorem [9] says that the local resilience of the complete graph K_n with respect to Hamiltonicity is at least $1/2 - o(1)$.

In recent years, an active and fruitful research direction in extremal and probabilistic combinatorics has become the study of resilience of random and pseudorandom structures. The systematic study of those with respect to various graph properties was initiated by Sudakov and Vu in [25], who in particular proved that $G(n, p)$ (i.e. the Erdős-Rényi random graph model that is defined on the vertex set $[n]$ with each pair of vertices forming an edge randomly and independently with probability p) has resilience at least $1/2 - o(1)$ with respect to Hamiltonicity a.a.s. for $p > \ln^4 n/n$. This result was improved by Lee and Sudakov [17] to $p \gg \ln n/n$, which is essentially best possible with respect to both the constant $1/2$ and the edge probability, since one can a.a.s. find disconnected spanning subgraphs of $G(n, p)$ with degree at most $(1/2 - o(1))pn$ and since $G(n, p)$ itself is a.a.s. disconnected for $p \leq (1 - o(1)) \ln n/n$.

A lot of resilience results are known for random graphs. For instance, the containment of triangle factors [4], almost spanning trees of bounded degree [3], pancyclic graphs [15], almost spanning and spanning bounded degree graphs with sublinear bandwidth [1, 6, 13], directed Hamilton cycles [11, 12, 18], perfect matchings and Hamilton cycles in random graph processes [19, 20], almost spanning powers of cycles [24] were studied.

An r -uniform hypergraph is a tuple (V, E) with $E \subseteq \binom{V}{r}$ and thus the generalisation of a graph: the elements of V are called *vertices* and the elements of E *hyperedges* (or *edges* for short). It is therefore natural to ask for degree conditions that force a subhypergraph (or *subgraph* for short) of the complete hypergraph to contain a copy of some given large structure. Such problems have been studied extensively in the last years, especially for different kinds of Hamilton cycles. Furthermore, (bounds on) the threshold for the existence of a Hamilton cycle in the random r -uniform hypergraph model $H^{(r)}(n, p)$ (every possible edge appears with probability p independently of the others) have been determined for various notions of cycles. We refer to [16] for an excellent survey by Kühn and Osthus of such problems.

The purpose of this work is to provide a Dirac-type result in the random hypergraph $H^{(r)}(n, p)$. This result was announced and its proof sketched in [7] and is, to the best of our knowledge, the first local resilience result in random hypergraphs at an almost optimal edge probability. The only other resilience result in random hypergraphs is a recent work of Ferber and Hirschfeld [10] on the resilience of perfect matchings with respect to the codegree condition in random hypergraphs at the asymptotically optimal probability.

Given an r -uniform hypergraph $H = (V, E)$. We use $\deg_H(v)$ to denote the vertex degree of a vertex v in H , i.e. the number of hyperedges of H that contain v . The minimum vertex degree of a hypergraph H is then $\delta_1(H) := \min_v \deg_H(v)$. We will also consider other degree notions such as $\deg_H(T)$ defined as $|\{e: e \supseteq T, e \in E(H)\}|$, i.e. the number of edges that contain a given tuple T . Generally, we define the maximum

$(\Delta_\ell(H))$ ℓ -collective degree as follows:

$$\Delta_\ell(H) := \max_{T \in \binom{V}{\ell}} \deg_H(T).$$

The notion of resilience in graphs extends verbatim to the setting of hypergraphs.

We will be interested in resilience results of random r -uniform hypergraphs with respect to weak and Berge Hamiltonicity. Weak and Berge cycles are the earliest notion of cycles in hypergraphs and are defined as follows.

Definition 1. A *weak cycle* is an alternating sequence $(v_1, e_1, v_2, \dots, v_k, e_k)$ of distinct vertices v_1, \dots, v_k and hyperedges e_1, \dots, e_k such that $\{v_1, v_k\} \subseteq e_k$ and $\{v_i, v_{i+1}\} \subseteq e_i$ for every $i \in [k-1]$. A weak cycle is called *Berge cycle* if all its hyperedges are distinct.

If $P = (v_1, e_1, v_2, \dots, v_n, e_n)$ is a weak cycle or a Berge cycle in a hypergraph H on n vertices, then P is called *weak Hamilton cycle* or *Hamilton Berge cycle* of H , respectively. Other common notion of cycles are ℓ -cycles. For an integer $1 \leq \ell \leq r-1$, an r -uniform hypergraph C is an ℓ -cycle if there exists a cyclic ordering of the vertices of C such that every hyperedge of C consists of r consecutive vertices and such that every pair of consecutive hyperedges intersects in precisely ℓ vertices.

If $\ell = 1$, then C is called a *loose cycle* and if $\ell = r-1$, then C is called a *tight cycle*. Observe that every ℓ -cycle is a Berge cycle. Furthermore, every tight Hamilton cycle (i.e. a spanning tight cycle) is a Hamilton Berge cycle. This is however not true for ℓ -cycles if $\ell < r-1$ since a Hamilton ℓ -cycle in a hypergraph on n vertices has $n/(r-\ell)$ hyperedges, whereas a Hamilton Berge cycle has n hyperedges. But since hyperedges may be repeated in weak cycles, a Hamilton ℓ -cycle is a weak Hamilton cycle. There is also an extensive literature on Dirac-type results that are stated in terms of the minimum degree for dense graphs and hypergraphs, we refer to the illuminating surveys [16, 22, 26].

Surprisingly, until recently, the only result on the minimum vertex degree which implies the existence of a weak or a Berge Hamilton cycle was the one due to Bermond, Germa, Heydemann, and Sotteau [5]. They proved that for every integer $r \geq 3$ and $k \geq r+1$ any r -uniform hypergraph H with minimum vertex degree $\delta_1(H) \geq \binom{k-2}{r-1} + r-1$ contains a Berge cycle on at least k vertices. If we ask for a Hamilton Berge cycle in an r -uniform hypergraph on n vertices, where r is fixed and n is large, then the bound $\binom{n-2}{r-1} + r-1$ is very weak since it differs from the maximum possible degree by $\binom{n-2}{r-2} - r+1$. Certainly, the two propositions below are folklore and should be known.

Proposition 2. Let $r \geq 3$ and $n \geq r$ and let H be an r -uniform hypergraph on n vertices. If $\delta_1(H) > \binom{\lceil n/2 \rceil - 1}{r-1}$, then H contains a weak Hamilton cycle.

The proof is a one-line argument by replacing every edge of H with a clique on r vertices and applying the original theorem by Dirac. The bound on the minimum vertex degree is sharp. Indeed, for even n , the disjoint union of two copies of the complete r -uniform hypergraph $K_{\frac{n}{2}}^{(r)}$ on $\frac{n}{2}$ vertices has minimum vertex degree $\binom{n/2-1}{r-1}$ but is disconnected. For odd n , the hypergraph H on n vertices that is the composition of two copies of $K_{\lceil \frac{n}{2} \rceil}^{(r)}$

that share one vertex satisfies $\delta_1(H) = \binom{\lceil n/2 \rceil - 1}{r-1}$ but does not contain a weak Hamilton cycle.

The following result can be obtained along the lines of the proof of Dirac's theorem for graphs.

Proposition 3. *Let $r \geq 3$ and let H be an r -uniform hypergraph on $n > 2r - 2$ vertices. If $\delta_1(H) \geq \binom{\lceil n/2 \rceil - 1}{r-1} + n - 1$ then H contains a Hamilton Berge cycle.*

In any case, it follows from Propositions 2 and 3 that the resilience of the complete hypergraph $K_n^{(r)}$ is at least $1 - 2^{1-r} - o(1)$ with respect to both weak and Berge Hamiltonicity. Recently, Coulson and Perarnau [8] improved the lower bound from Proposition 3 for sufficiently large n to the optimal $\delta_1(H) \geq \binom{\lceil n/2 \rceil - 1}{r-1}$.

Like in the setting of graphs, a natural question is which sparse random hypergraphs contain a weak Hamilton cycle or even a Hamilton Berge cycle and how robust these hypergraphs are with respect to these properties? Recall that by $H^{(r)}(n, p)$ we denote the random r -uniform hypergraph model on the vertex set $[n]$, where each set of r vertices forms an edge randomly and independently with probability $p = p(n)$. Poole [21] determined the threshold for the existence of a weak Hamilton cycle in $H^{(r)}(n, p)$.

Theorem 4 (Theorem 1.1 in [21]). *Let $r \geq 3$. Then*

$$\mathbb{P}[H^{(r)}(n, p) \text{ is weak Hamiltonian}] \rightarrow \begin{cases} 0 & \text{if } p \leq (r-1)! \frac{\ln n - \omega(1)}{n^{r-1}} \\ e^{-e^{-c}} & \text{if } p = (r-1)! \frac{\ln n + c_n}{n^{r-1}} \\ 1 & \text{if } p \geq (r-1)! \frac{\ln n + \omega(1)}{n^{r-1}}, \end{cases}$$

where c_n is any function tending to $c \in \mathbb{R}$ and $\omega(1)$ is an arbitrary function that tends to infinity.

Since every Hamilton Berge cycle is in particular a weak Hamilton cycle, Theorem 4 yields that $H^{(r)}(n, p)$ a.a.s. does not contain a Hamilton Berge cycle if $p \leq (r-1)! \frac{\ln n - \omega(1)}{n^{r-1}}$. Surprisingly, before our extended abstract [7] an upper bound on the threshold of $H^{(r)}(n, p)$ being Berge Hamiltonian has not been studied directly, but recently, Bal, Berkowitz, Devlin and Schacht [2] determined the threshold for Hamilton Berge cycles up to the constant factor.

Main result

In this paper we prove the following Dirac-type result for the existence of Hamilton Berge cycles in random hypergraphs.

Theorem 5. *For every integer $r \geq 3$ and every real $\gamma > 0$ the following holds asymptotically almost surely for $\mathcal{H} = H^{(r)}(n, p)$ if $p \geq \frac{\log^{17r} n}{n^{r-1}}$. Let $H \subseteq \mathcal{H}$ be a spanning subgraph with $\delta_1(H) \geq \left(\frac{1}{2^{r-1}} + \gamma\right) p \binom{n}{r-1}$. Then H contains a Hamilton Berge cycle.*

It is worth noting that it is not possible to merely reduce the problem of finding Berge cycles to a problem of finding a Hamilton cycle in the random graph to achieve the same resilience $1 - 2^{1-r} + o(1)$ in the random hypergraph $H^{(r)}(n, p)$. One possibility of such a ‘reduction’ would be to declare an edge $uv \in \binom{[n]}{2}$ to lie in G if $\{u, v\}$ lies in some hyperedge of $H^{(r)}(n, p)$. Then such G behaves much like the random graph $G(n, p')$, where $1 - p' = (1 - p)^{\binom{n-2}{r-2}}$, i.e. $p' = \Omega(\text{polylog}(n)/n)$. However, such a reduction leads to a resilience which is far from being the asymptotically optimal one, as asserted by our Theorem 5 (lower bound) and an asymptotically matching upper bound discussed below.

The minimum degree condition is asymptotically tight, meaning that we cannot replace it with $\delta_1(H) \geq (\frac{1}{2^{r-1}} - \gamma) p \binom{n}{r-1}$. Indeed, given $\mathcal{H} = H^{(r)}(n, p)$ together with a partition $V(\mathcal{H}) = V_1 \cup V_2$ with $|V_1 - V_2| \leq 1$, it happens a.a.s. that, for $i = 1, 2$, the degree of every $v \in V_i$ into V_i is at least $(\frac{1}{2^{r-1}} - \gamma) p \binom{n}{r-1}$. The hypergraph $H := \mathcal{H}[V_1] \cup \mathcal{H}[V_2]$ does not contain a Berge Hamilton cycle and satisfies $\delta_1(H) \geq (\frac{1}{2^{r-1}} - \gamma) p \binom{n}{r-1}$.

Furthermore, the bound on p is optimal up to possibly this polylogarithmic factor, and it provides an alternative proof of the result in [21] with only slightly weaker edge probability. The proof is based on the absorbing method developed by Rödl, Ruciński, and Szemerédi [23]. Of particular importance are the ideas from the proof of a Dirac-type result for random directed graphs due to Ferber, Nenadov, Noever, Peter and Škoric [11], which allow us to apply this method in such a very sparse scenario.

Organization of the paper

The rest of the paper is organised as follows. In Section 2 we state concentration inequalities that we are going to use. In Section 3 we introduce some technical definitions, the notion of (ε, p) -pseudorandom hypergraphs and prove some basic properties around these notions. In Section 4 we introduce the central concept of absorbers and in Section 5 we prove a connection lemma which will allow us to connect many pairs of vertices by disjoint paths. We provide in Section 6 first an outline of the proof of Theorem 5 and then we prove a technical theorem about Hamilton Berge cycles in pseudorandom hypergraphs which implies Theorem 5. We present the proof of Proposition 3 in Section 7.

2 Probabilistic tools: concentration inequalities

In our proofs we use the following standard bounds on deviations of random variables. Their proofs can be found in e.g. [14]. The following two bounds belong to Chernoff’s inequality, which collects different exponentially decreasing bounds on the tails of a binomial distribution.

Theorem 6 (Chernoff’s inequality I). *For every random variable $X \sim \text{Bin}(n, p)$ and every $\varepsilon \leq 3/2$ we have*

$$\mathbb{P}[|X - \mathbb{E}[X]| > \varepsilon \mathbb{E}[X]] < 2 \exp\left(-\frac{\varepsilon^2 \mathbb{E}[X]}{3}\right).$$

The second Chernoff's inequality that we need provides only a bound on the upper tail of the binomial distribution.

Theorem 7 (Chernoff's inequality II). *For every random variable $X \sim \text{Bin}(n, p)$ and every $t \geq 0$ we have*

$$\mathbb{P}[X \geq \mathbb{E}[X] + t] \leq \exp\left(-\frac{t^2}{2(\mathbb{E}[X] + t/3)}\right).$$

Finally we consider binomial random subsets. For $\Gamma = [n]$ let Γ_{p_1, \dots, p_n} be defined by including for every $i \in [n]$ the i -th element of Γ with probability p_i independently of all other elements of Γ . For each set $\mathcal{S} \subseteq 2^\Gamma$ of subsets of Γ and each set $A \in \mathcal{S}$, we let X_A denote the indicator variable for the event $A \subseteq \Gamma_{p_1, \dots, p_n}$. Janson's inequality gives an exponentially small bound on the lower tail of the distribution of sums of such indicator variables.

Theorem 8 (Janson's inequality). *Let Γ be a finite set and let $\mathcal{S} \subseteq 2^\Gamma$ be a set of subsets of Γ . If $X = \sum_{A \in \mathcal{S}} X_A$, where X_A is an indicator variable, and $0 \leq t \leq \mathbb{E}[X]$, then*

$$\mathbb{P}[X \leq \mathbb{E}[X] - t] \leq \exp\left(-\frac{t^2}{2\overline{\Delta}}\right),$$

where

$$\overline{\Delta} = \mathbb{E}[X] + \sum_{A \in \mathcal{S}} \sum_{\substack{B \in \mathcal{S}: \\ A \cap B \neq \emptyset, A \neq B}} \mathbb{E}[X_A X_B].$$

3 More definitions and auxiliary lemmas

Given an r -uniform hypergraph H , a subset $U \subseteq V(H)$ and a vertex $u \in V(H)$, we write $\deg_H(u, U)$ to denote the degree of u in U , which is the number of edges from H which contain u and which are contained completely in $U \cup \{u\}$. For disjoint T and U , we write $e_H(T, \binom{U}{r-1})$ to denote the number of edges e from H with $|e \cap T| = 1$ and $|e \cap U| = r - 1$. The hypergraph H is sometimes omitted when it is clear from the context. By an ℓ -set we mean an ℓ -element set and by an ℓ -subset we refer to an ℓ -element subset of some set.

Our proof will use some central properties of the random hypergraph $H^{(r)}(n, p)$ which we call pseudorandom. So, our main theorem will state that r -uniform hypergraphs that are pseudorandom satisfy a Dirac-type theorem about Hamilton Berge cycles.

3.1 Pseudorandom hypergraphs

The following proposition asserts that for an edge probability $p \leq \text{polylog}(n)/n^{r-1}$, the 2-collective degree of a random hypergraph does not get too large.

Proposition 9. *For every integer $r \geq 3$, every real $c > 0$ and $p \leq \frac{\ln^c n}{n^{r-1}}$, with probability at least $1 - 1/n$, the hypergraph $H^{(r)}(n, p)$ has maximum 2-collective degree at most $2 \ln n$.*

Proof. The statement is an immediate consequence of Chernoff's inequality, Theorem 7, and the union bound over $\binom{n}{2}$ possible pairs of vertices. \square

Next we verify that the edges in the random hypergraph $H^{(r)}(n, p)$ are distributed as expected.

Lemma 10. *For every integer $r \geq 3$, reals $\varepsilon, p > 0$ and sufficiently large integer n the following holds with probability at least $1 - \frac{2}{n}$ for any pair of disjoint subsets T and U of $[n]$ with $|U| \leq |T|$:*

$$e_{H^{(r)}(n,p)}\left(T, \binom{U}{r-1}\right) \leq (1 + \varepsilon)p|T|\binom{|U|}{r-1} + |T| \ln^{1+\varepsilon} n.$$

Proof. Consider the random variable $X = e_{H^{(r)}(n,p)}\left(T, \binom{U}{r-1}\right)$, then its expectation is clearly $\mathbb{E}[X] = p|T|\binom{|U|}{r-1}$. If $\mathbb{E}[X] \geq \frac{12}{\varepsilon^2}|T| \ln n$, then Chernoff's inequality, Theorem 7, with $t = \varepsilon\mathbb{E}[X]$ gives

$$\mathbb{P}[X \geq (1 + \varepsilon)\mathbb{E}[X]] \leq \exp\left(-\frac{\varepsilon^2}{3}\mathbb{E}[X]\right) \leq n^{-4|T|}.$$

If $\mathbb{E}[X] < \frac{12}{\varepsilon^2}|T| \ln n$, then Chernoff's inequality, Theorem 7, with $t = |T| \ln^{1+\varepsilon} n$ gives

$$\mathbb{P}[X \geq \mathbb{E}[X] + t] \leq \exp(-|T| \ln^{1+\varepsilon} n) \leq n^{-4|T|}.$$

Since there are at most $\binom{n}{|T|}\binom{n}{|U|} \leq n^{2|T|}$ pairs (T, U) with fixed sizes satisfying $|T| \geq |U|$ we use the union bound to obtain that with probability at least $1 - n(\sum_{i=1}^n n^{2i} n^{-4i}) \geq 1 - \frac{2}{n}$ for all pairs of disjoint sets (T, U) with $|T| \geq |U|$ it holds that

$$e_{H^{(r)}(n,p)}\left(T, \binom{U}{r-1}\right) \leq (1 + \varepsilon)p|T|\binom{|U|}{r-1} + |T| \ln^{1+\varepsilon} n. \quad \square$$

Lemma 11. *For every integer $r \geq 3$, reals $\varepsilon, p > 0$ with $\varepsilon \in (0, 3/2)$ and sufficiently large integer n , the following holds with probability at least $1 - \frac{2}{n}$ for any pair of disjoint subsets T and U with $\varepsilon|U| \leq |T| \leq |U| \leq n/2$ and $|U| \geq \left(\frac{13(r-1)!\ln n}{\varepsilon^3 p}\right)^{1/(r-1)}$:*

$$e_{H^{(r)}(n,p)}\left(T, \binom{U}{r-1}\right) \leq (1 + \varepsilon)p|T|\binom{|U|}{r-1}.$$

Proof. Consider the random variable $X = e_{H^{(r)}(n,p)}\left(T, \binom{U}{r-1}\right)$, then its expectation is clearly $\mathbb{E}[X] = p|T|\binom{|U|}{r-1}$. Set $m := \left(\frac{13(r-1)!\ln n}{\varepsilon^3 p}\right)^{1/(r-1)}$. By Chernoff's inequality, Theorem 6, we have for fixed disjoint sets U and T with $|T| \geq \varepsilon|U| \geq \varepsilon m$

$$\mathbb{P}[X \geq (1 + \varepsilon)\mathbb{E}[X]] \leq 2 \exp\left(-\frac{\varepsilon^2\mathbb{E}[X]}{3}\right) \leq 2 \exp(-\varepsilon^3 p|U|^r / (4(r-1)!)) \leq n^{-3|U|}.$$

Since there are at most $\binom{n}{|T|}\binom{n}{|U|} \leq (en/|U|)^{2|U|}$ pairs (T, U) of fixed sizes $|T|$ and $|U|$ we use the union bound to obtain that with probability at least $1 - \sum_{i=m}^{n/2} \sum_{j=\varepsilon i}^i (en/i)^{2i} n^{-3i} \geq 1 - \frac{2}{n}$ for all pairs of disjoint sets (T, U) it holds that

$$e_{H^{(r)}(n,p)}\left(T, \binom{U}{r-1}\right) \leq (1 + \varepsilon)p|T|\binom{|U|}{r-1}. \quad \square$$

Lemmas 10 and 11 motivate the following definition of an (ε, p) -pseudorandom hypergraph.

Definition 12 ((ε, p) -pseudorandomness). Given $p > 0$ and $\varepsilon \in (0, 3/2)$, a hypergraph \mathcal{H} on n vertices with $\Delta_2(\mathcal{H}) \leq 2 \ln n$ is (ε, p) -pseudorandom if it satisfies the following properties:

- (i) for any pair of disjoint subsets T and U of $[n]$ with $|U| \leq |T|$

$$e_{\mathcal{H}}\left(T, \binom{U}{r-1}\right) \leq (1 + \varepsilon)p|T|\binom{|U|}{r-1} + |T| \ln^{1+\varepsilon} n,$$

- (ii) for any pair of disjoint subsets T and U with $\varepsilon|U| \leq |T| \leq |U| \leq n/2$ and $|U| \geq \left(\frac{13(r-1)! \ln n}{\varepsilon^3 p}\right)^{1/(r-1)}$

$$e_{H^{(r)}(n,p)}\left(T, \binom{U}{r-1}\right) \leq (1 + \varepsilon)p|T|\binom{|U|}{r-1}.$$

3.2 A sampling lemma

We will repeatedly use the fact that in a hypergraph of high minimum degree, a random subset of vertices inherits high minimum degree of every vertex. More precisely, we prove the following.

Lemma 13. Let $r \geq 3$, reals $\gamma > 0$, $c' > 0$ and $c > 2 + 2c'$ be given. Let $H = (V, E)$ be an r -uniform hypergraph on n vertices and V' be a subset of V with at least $2m$ vertices, where $m \geq \frac{n}{\ln^c n}$. Furthermore assume that $\deg_H(v, V') \geq \gamma p \binom{|V'|}{r-1}$ for all $v \in V$ and $\Delta_2(H) \leq 2 \ln n$, where $p \geq \frac{\ln^{cr} n}{n^{r-1}}$. Then the following holds for all n sufficiently large.

There exists a set $U \subset V'$ of size m such that

$$(i) \deg_H(v, U) \geq \left(1 - \frac{1}{\ln^{c'} n}\right) \gamma p \binom{|U|}{r-1} \text{ for all } v \in V, \text{ and}$$

$$(ii) \deg_H(v, V' \setminus U) \geq \left(1 - \frac{1}{\ln^{c'} n}\right) \gamma p \binom{|V'| - |U|}{r-1} \text{ for all } v \in V.$$

Proof. We choose a set U randomly, by including every vertex u from V' into U with probability $q = \frac{m}{n'}$ independently, where we set $n' := |V'|$. Thus, we associate with every vertex u a Bernoulli variable t_u with parameter q .

For a given vertex $v \in V$ let X_v be the random variable for $\deg_H(v, U)$. Clearly, we can write X_v as the following sum of indicator random variables:

$$X_v = \sum_{A \in \mathcal{S}} X_A,$$

where $X_A = \prod_{u \in A} t_u$ and $\mathcal{S} = \{e \setminus \{v\} : e \in E(H), v \in e\}$.

We are going to apply Theorem 8 and for this we put the following estimates: $\mathbb{E}[X_v] = q^{r-1} \deg_H(v) \geq q^{r-1} \gamma p \binom{n'}{r-1} \geq \gamma \ln^c n / (2(r-1)!)$, and

$$\begin{aligned} \overline{\Delta} &= \mathbb{E}[X_v] + \sum_{A \in \mathcal{S}} \sum_{\substack{B \in \mathcal{S}: \\ A \cap B \neq \emptyset, A \neq B}} \mathbb{E}[X_A X_B] \\ &\leq \mathbb{E}[X_v] + \deg_H(v)(r-1) \Delta_2(H) q^r \leq \mathbb{E}[X_v](1 + 2rq \ln n). \end{aligned}$$

Then Theorem 8 yields

$$\mathbb{P} \left[X_v \leq \mathbb{E}[X_v] - \frac{1}{\ln^{c'} n} \mathbb{E}[X_v] \right] \leq \exp \left(- \frac{\mathbb{E}[X_v]^2}{2(\ln^{2c'} n) \mathbb{E}[X_v](1 + 2rq \ln n)} \right) \leq \exp \left(- \frac{\mathbb{E}[X_v]}{4r \ln^{1+2c'} n} \right) < n^{-3}.$$

A similar argument applies also to $V' \setminus U$. Taking the union bound, with probability, say, at least $1 - \frac{1}{n}$ the properties (i) and (ii) are satisfied by all vertices $v \in V$.

On the other hand

$$\begin{aligned} \mathbb{P}[|U| = m] &= \mathbb{P}[|V' \setminus U| = n' - m] = \binom{n'}{m} q^m (1 - q)^{n' - m} = \frac{(1 + o(1))}{\sqrt{2\pi q(n' - m)}} \\ &\geq (1 + o(1)) \sqrt{\frac{2}{\pi n'}}. \end{aligned}$$

Therefore, with positive probability (at least $(1 + o(1)) \sqrt{\frac{2}{\pi n'}} - \frac{1}{n}$) there exists a desired set U . \square

3.3 Matchings

Our building blocks for Hamilton cycles will consist of collections of edges between pairs of equal-sized sets, which we will refer to as (U_1, U_2) -matchings. Moreover, these edges will intersect both sets U_1 and U_2 in a clearly specified way. The following two definitions make these ideas precise.

Definition 14 (An (i, j) -edge for (U_1, U_2)). Given a uniform hypergraph $H = (V, E)$, two disjoint subsets $U_1, U_2 \subseteq V$ and an edge $e \in E$. We call e an (i, j) -edge for (U_1, U_2) , if $|e \cap U_1| = i$ and $|e \cap U_2| = j$ hold.

Definition 15 ((U_1, U_2) -matching in H). Given an r -uniform hypergraph $H = (V, E)$ and two disjoint subsets $U_1, U_2 \subseteq V$ with $|U_2| \geq |U_1| = m$. We call the set $M = \{e_1, \dots, e_m\}$ a (U_1, U_2) -matching in H if there exists a matching $M' = \{a_i b_i : a_i \in U_1, b_i \in U_2, i \in [m]\}$ in the complete bipartite graph K_{U_1, U_2} with classes U_1 and U_2 such that $a_i, b_i \in e_i$ and e_i is a $(1, r-1)$ -edge or an $(r-1, 1)$ -edge for (U_1, U_2) for every $i \in [m]$.

We call the vertices a_i and b_i the *endpoints of the matching edge* e_i .

The next lemma asserts that between two disjoint subsets U_1, U_2 of vertices of high ‘minimum degree’ and of size $m = n/\text{polylog}(n)$ there must always be a (U_1, U_2) -matching M which intersects only U_1 and U_2 in the ‘pattern’ $(1, r-1)$ or $(r-1, 1)$. Its proof is an application of Hall’s matching criterion under the exploitation of the properties of (ε, p) -pseudorandom hypergraphs.

Lemma 16. *For every integer $r \geq 3$, every real $\gamma \in (0, 1)$ and $c > 1$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} on n vertices with n sufficiently large and $p \geq \frac{\ln^{cr} n}{n^{r-1}}$. Let $H \subseteq \mathcal{H}$ be a subgraph of \mathcal{H} and let U_1, U_2 be disjoint subsets of $V(H)$ with $|U_1| = |U_2| = m \geq \frac{n}{\ln^c n}$ such that $\deg_H(u, U_i) \geq \left(\frac{1}{2^{r-1}} + \gamma\right) p \binom{|U_i|}{r-1}$ for every $u \in U_{3-i}$ and $i = 1, 2$. Then there exists a (U_1, U_2) -matching in H .*

Proof. We choose with foresight $\varepsilon \leq 2^{r-3}\gamma$ such that $1 + \varepsilon < c$ holds.

For $T \subseteq U_1$ we define the neighbourhood $N(T) \subseteq U_2$ as follows (and similarly one defines $N(T) \subseteq U_1$ for $T \subseteq U_2$)

$$N(T) = \{b : \exists a \in T, e \in H \text{ with } a, b \in e, |e \cap U_2| = r-1\}.$$

It will be sufficient to verify $|N(T)| \geq |T|$ for sets $T \subseteq U_i$ (Hall’s condition) with $|T| \leq \lceil m/2 \rceil$ where $i \in [2]$. Assume w.l.o.g. that $T \subseteq U_1$ and further suppose towards a contradiction that $|N(T)| < |T| \leq \lceil m/2 \rceil$. Then, by the assumptions of the lemma, we have

$$e_H \left(T, \binom{U_2}{r-1} \right) = \sum_{u \in T} \deg_H(u, U_2) \geq |T| \left(\frac{1}{2^{r-1}} + \gamma \right) p \binom{m}{r-1}.$$

On the other hand, it follows from the (ε, p) -pseudorandomness property (i) of \mathcal{H} , that

$$e_H \left(T, \binom{U_2}{r-1} \right) = e_H \left(T, \binom{N(T)}{r-1} \right) \leq (1 + \varepsilon) p |T| \binom{|N(T)|}{r-1} + |T| \ln^{1+\varepsilon} n.$$

We estimate further $(1 + \varepsilon) p |T| \binom{|N(T)|}{r-1} < (1 + \varepsilon) p |T| \binom{\lceil m/2 \rceil}{r-1} \leq \frac{1+2\varepsilon}{2^{r-1}} p |T| \binom{m}{r-1}$. Comparing $\varepsilon 2^{r-2} p |T| \binom{m}{r-1} + |T| \ln^{1+\varepsilon} n$ with $\gamma p |T| \binom{m}{r-1}$, we obtain a contradiction in view of the choice of ε and the estimate $p \binom{m}{r-1} > \frac{\ln^{cr} n}{n^{r-1}} \frac{m^{r-1}}{r^{r-1}} \geq \frac{\ln^c n}{n^{r-1}}$. \square

Analogously to Lemma 16, one may prove almost verbatim the next lemma which asserts that between two disjoint subsets U_1, U_2 with $|U_2|/2 \geq |U_1| = m = n/\text{polylog}(n)$ and with all vertices in U_1 having high degree into U_2 there must always be a (U_1, U_2) -matching M consisting only of $(1, r-1)$ -edges for (U_1, U_2) .

Lemma 17. *For every integer $r \geq 3$, every real $\gamma \in (0, 1)$ and $c \geq 2$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} on n vertices with n sufficiently large and $p \geq \frac{\ln^{cr} n}{n^{r-1}}$. Let $H \subseteq \mathcal{H}$ be a subgraph of \mathcal{H} and let U_1, U_2 be disjoint subsets of $V(H)$ with $|U_2|/2 \geq |U_1| \geq \frac{n}{\ln^c n}$ such that $\deg_H(u, U_2) \geq \left(\frac{1}{2^{r-1}} + \gamma\right) p \binom{|U_2|}{r-1}$ for every $u \in U_1$. Then there exists a (U_1, U_2) -matching in H consisting only of $(1, r-1)$ -edges for (U_1, U_2) . \square*

The next technical definition is very handy to describe the basic structures we will be interested in.

Definition 18. Given an r -uniform hypergraph H and two disjoint sets A and B with $|B| \geq 2|A|$. A 2-matching for (A, B) is a collection of pairs of edges $(e_a, f_a)_{a \in A}$ so that

- (i) all these edges are distinct,
- (ii) $a \in e_a, a \in f_a$ for all $a \in A$,
- (iii) there is an injection $\tau: \cup_{a \in A} \{e_a, f_a\} \rightarrow B$ with $\tau(g) \in g$, and
- (iv) for every edge $g \in \cup_{a \in A} \{e_a, f_a\}$: $|g \cap A| = 1$ and $|g \cap B| = r - 1$.

The next lemma allows us to find a 2-matching.

Lemma 19. For every integer $r \geq 3$, every real $\gamma \in (0, 2^{1-r})$ and $c > 2$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} on n vertices with $p \geq \frac{\ln^{cr} n}{n^{r-1}}$ and n sufficiently large with $\Delta_2(\mathcal{H}) \leq 2 \ln n$. Let $H \subseteq \mathcal{H}$ be a subgraph of \mathcal{H} and let A and B be disjoint subsets of $V(H)$ with $|A| = m$ and $|B| \geq 4m$, where $m \geq \frac{n}{\ln^c n}$, and such that $\deg_H(a, B) \geq (\frac{1}{2^{r-1}} + 2.5\gamma) p \binom{|B|}{r-1}$ for every $a \in A$. Then there exists a 2-matching for (A, B) .

Proof. We apply Lemma 13 to B and obtain an equipartition into $B_1 \cup B_2$ such that $\deg_H(a, B_i) \geq (\frac{1}{2^{r-1}} + 2\gamma) p \binom{|B_i|}{r-1}$ for every $a \in A$ and $i \in [2]$. An application of Lemma 17 to (A, B_1) and (A, B_2) yields the desired 2-matching. \square

4 Absorbers

A *weak Berge path* (or simply *weak path*) is an alternating sequence $(v_1, e_1, v_2, \dots, v_k)$ of distinct vertices v_1, \dots, v_k and (not necessarily distinct) hyperedges e_1, \dots, e_{k-1} such that $v_i, v_{i+1} \in e_i$ for every $i \in [k-1]$. A weak path is called *Berge path* if all its hyperedges are distinct.

For a weak path $P = (v_1, e_1, \dots, e_{k-1}, v_k)$ we denote by $E(P) := \{e_1, \dots, e_{k-1}\}$ the set of hyperedges of P , by $V^*(P) := \{v_1, \dots, v_k\}$ the set of *inner* vertices in the sequence of P , and by $V(P) := \bigcup_{i \in [k-1]} e_i$ the union of the hyperedges of P . We say that P *connects* v_1 to v_k and call v_1 and v_k *endpoints* of P . For a weak or Berge cycle $C = (v_1, e_1, v_2, \dots, v_k, e_k)$ we define $V^*(C) := \{v_1, \dots, v_k\}$ and we refer to (v_1, \dots, v_k) as a *sequence* of C .

The *length* of a weak path P is defined as $|V^*(P)| - 1$, and the *length* of a weak cycle C is defined as $|V^*(C)|$. In particular, if P is a Berge path, then the length of P is exactly the number of hyperedges of P . Given two weak paths $P = (v_1, e_1, \dots, e_{k-1}, v_k)$ and $Q = (v_k, e'_1, \dots, e'_{k'-1}, v'_{k'})$ with $V^*(P) \cap V^*(Q) = \{v_k\}$, we denote by $P \cdot Q$ the weak path $(v_1, e_1, \dots, e_{k-1}, v_k, e'_1, \dots, e'_{k'-1}, v'_{k'})$.

We say that two Berge paths $P = (v_1, e_1, \dots, e_{k-1}, v_k)$ and $P' = (v'_1, e'_1, \dots, e'_{k'-1}, v'_{k'})$ are *edge-disjoint* if $e_i \neq e'_j$ for all $i \in [k-1]$ and $j \in [k'-1]$.

Next we introduce the notion of an absorber.

Definition 20 (Absorber for a vertex u). For a uniform hypergraph H and a vertex u , an absorber for u is a subgraph A of H which consists of the following edges specified in the properties below:

- (i) A contains a Berge cycle C with $u \in V^*(C)$ of length $2t + 1$ for some $t \in \mathbb{N}$ and with vertex sequence $(u, v_1, \dots, v_{t+1}, \dots, v_{2t})$;
- (ii) there are $t - 1$ Berge paths P_1, \dots, P_{t-1} so that each path P_i has endpoints v_{i+1} and v_{2t+1-i} and the inner vertex sets are pairwise disjoint;
- (iii) the edge-sets $E(C), E(P_1), \dots, E(P_{t-1})$ are pairwise disjoint;
- (iv) $E(A) = E(C) \cup \bigcup_{i \in [t-1]} E(P_i)$.

We call the vertex u a *reservoir vertex* and the absorber A a u -*absorber*. The *inner vertices* of A are the vertices from $V^*(C) \cup \bigcup_i V^*(P_i)$. The vertices v_1 and v_{t+1} are referred to as the *main endpoints* of A .

The following proposition about an absorber for some vertex u explains its extreme usefulness in what comes and also the role of the main endpoints of an absorber – the absorber A_u contains two Berge paths with the same endpoints and the inner vertices of the first path consist of all inner vertices of the absorber, while the inner vertices of the second contain all inner vertices but u .

Proposition 21. *Let A be a u -absorber in some hypergraph H and let C and P_1, \dots, P_{t-1} be the Berge cycle and Berge paths of the absorber A respectively. Let $(u, v_1, \dots, v_{t+1}, \dots, v_{2t})$ be the sequence of C according to property (i) in Definition 20. Then A contains the following two Berge paths:*

- (a) a Berge path P_u from v_1 to v_{t+1} with $V^*(P_u) = V^*(C) \cup \bigcup_{i \in [t-1]} V^*(P_i)$, and
- (b) a Berge path P from v_1 to v_{t+1} with $V^*(P) = (V^*(C) \setminus \{u\}) \cup \bigcup_{i \in [t-1]} V^*(P_i)$.

Proof. Let the structure of the cycle C be as follows

$$C = (u, e_1, v_1, e_2, \dots, e_{t+1}, v_{t+1}, \dots, v_{2t}, e_{2t+1}).$$

W.l.o.g. assume t is even (the case t odd is very similar). Now P_u is given as

$$P_u = (v_1, e_1, u, e_{2t+1}, v_{2t}) \cdot P_1 \cdot (v_2, e_3, v_3) \cdot P_2 \cdot \dots \cdot P_{t-1} \cdot (v_t, e_{t+1}, v_{t+1}).$$

Then the path P is defined as follows

$$P = (v_1, e_2, v_2) \cdot P_1 \cdot (v_{2t}, e_{2t-1}, v_{2t-1}) \cdot P_2 \cdot \dots \cdot P_{t-1} \cdot (v_{t+2}, e_{t+1}, v_{t+1}).$$

It is most instructive to draw a picture: placing the inner vertices consecutively on a cycle and connecting appropriate vertices with the paths, one sees immediately that there is one way to traverse all vertices and a ‘complementary’ way to traverse all vertices except for u . Since the paths P_i and the cycle C in the absorber use different edges, it follows that these weak paths are indeed Berge. \square

5 Connection lemma

In this section we will concentrate on a connection lemma that will allow us to put Berge paths together into a longer Berge path.

5.1 An expansion lemma

The following lemma allows us to prove an expansion property for a pseudorandom hypergraph in a resilience setting between any two ‘random’, not too small vertex subsets.

Lemma 22. *For every integer $r \geq 3$, every real $\gamma \in (0, 2^{1-r})$ and $c > 1$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} with $p \geq \frac{\ln^{cr} n}{n^{r-1}}$. Let $H \subseteq \mathcal{H}$ be a subgraph of \mathcal{H} and let U_1, U_2 be disjoint subsets of $V(H)$ with $|U_1| = |U_2| = m \geq \frac{n}{\ln^c n}$ such that $\deg_H(u, U_2) \geq \left(\frac{1}{2^{r-1}} + 2\gamma\right) p \binom{m}{r-1}$ for every $u \in U_1$. Then for every subset $T_1 \subseteq U_1$ of cardinality at least γm there exists a subset $T_2 \subseteq U_2$ of cardinality at least $(1/2 + \gamma)m$ such that for every $b \in T_2$ there exists a $(1, r-1)$ -edge e for (T_1, T_2) with $b \in e$.*

Proof. We choose $\varepsilon > 0$ such that $(1 + \varepsilon)\left(\frac{1}{2} + \gamma\right)^{r-1} < 2^{1-r} + 2\gamma$.

Let an (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} and a subhypergraph $H \subseteq \mathcal{H}$ with sets U_1, U_2 as specified in the assumption of the lemma be given. Without loss of generality let $T_1 \subseteq U_1$ with $|T_1| = \gamma m$. We define $T_2 := \{b : \exists a \in T_1, e \in H \text{ with } a, b \in e, |e \cap U_2| = r-1\}$ and assume that $|T_2| < (1/2 + \gamma)m$. First we arbitrarily extend T_2 to a subset $T'_2 \subseteq U_2$ of size $(1/2 + \gamma)m$.

Then compare the lower bound $e_H(T_1, \binom{U_2}{r-1}) = e_H(T_1, \binom{T'_2}{r-1}) \geq |T_1| \left(\frac{1}{2^{r-1}} + 2\gamma\right) p \binom{m}{r-1}$ with the upper bound on $e_H(T_1, \binom{T'_2}{r-1})$ which comes from condition (ii) of the definition of (ε, p) -pseudorandomness (it is easily seen that $|T'_2| \geq \left(\frac{13(r-1)! \ln n}{\varepsilon^3 p}\right)^{1/(r-1)}$ holds for n large enough):

$$\begin{aligned} |T_1| \left(\frac{1}{2^{r-1}} + 2\gamma\right) p \binom{m}{r-1} &\leq e_H \left(T_1, \binom{T_2}{r-1}\right) \leq e_H \left(T_1, \binom{T'_2}{r-1}\right) \\ &\leq (1 + \varepsilon) p \gamma m \binom{(1/2 + \gamma)m}{r-1}. \end{aligned}$$

Since $(1 + \varepsilon) p \gamma m \binom{(1/2 + \gamma)m}{r-1} < (1 + \varepsilon)(1/2 + \gamma)^{r-1} p \gamma m \binom{m}{r-1}$, we however obtain a contradiction due to our choice of ε . Thus, T_2 must contain more than $(1/2 + \gamma)m$ vertices. \square

5.2 A connection lemma

Our goal is to prove a connection lemma, Lemma 25, which will allow us to connect a given collection of pairs of vertices by edge-disjoint Berge paths, all of whose inner vertices but possibly endpoints are pairwise disjoint as well. The lemma is similar to the connection lemma from [11].

First we argue that in a subgraph of a pseudorandom hypergraph there exists a Berge path between a sequence of disjoint, not so small sets, even after deleting a positive proportion of the vertices from each of the sets, if one has sufficiently high minimum vertex degree between every two ‘consecutive’ sets.

Lemma 23. *For every integer $r \geq 3$, every real $\gamma \in (0, 2^{1-r})$ and $c > 1$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} on n vertices with $p \geq \frac{\ln^{cr} n}{n^{r-1}}$ and n sufficiently large. Let $H \subseteq \mathcal{H}$ be a subgraph of \mathcal{H} , let $m \geq \frac{n}{\ln^c n}$ and U_1, U_2, \dots, U_k (for some $k \geq \log_2 m$) be pairwise disjoint subsets of $V(H)$ each of cardinality m such that $\deg_H(u, U_{i+1}) \geq (\frac{1}{2^{r-1}} + 2\gamma) p \binom{m}{r-1}$ for every $u \in U_i$ and every $i \in [k-1]$. Then for every sequence W_1, \dots, W_{k-1} of subsets $W_i \subseteq U_i$, $i \in [k-1]$, such that $|W_1| = \gamma m$ and $|W_i| \geq (1 - \gamma)m$ for $i \geq 2$, there exists a vertex $v_1 \in W_1$ and a set $T_k \subseteq U_k$ of cardinality at least $(1/2 + \gamma)m$ with the following property:*

- *for every $v_k \in T_k$ there is a Berge path $P = (v_1, e_1, \dots, e_{k-1}, v_k)$ with $v_i \in W_i$ for all $i \in [k-1]$ and such that $|e_i \cap U_{i+1}| = r-1$ for all $i \in [k-1]$.*

Proof. We choose $\varepsilon > 0$ so that Lemma 22 is applicable on input r, γ and c .

First we show the following statement for all $j = 2, \dots, k$:

- (\star) there exists a subset $T_1 \subseteq W_1$ of cardinality $\max\{1, \gamma m/2^{j-1}\}$ and a subset $T_j \subseteq U_j$ of cardinality at least $(1/2 + \gamma)m$ such that for every $v_j \in T_j$ there is a Berge path $P = (v_1, e_1, \dots, e_{j-1}, v_j)$ with $v_i \in W_i$ for all $i \in [j-1]$ and such that $|e_i \cap U_{i+1}| = r-1$ for all $i \in [j-1]$.

The case $j = 2$ follows from Lemma 22 (all of whose assumptions are met): the Berge paths correspond to single edges with appropriate endpoints.

Let now $j > 2$ and proceed by induction. Then, we use the truth of statement (\star) for $j-1$ and let T_1 and T_{j-1} be the corresponding sets. We have $|T_{j-1} \cap W_{j-1}| \geq m/2$. We partition T_1 into two equal-sized sets T'_1 and T''_1 (if $|T_1|=1$ then we set $T'_1 := T''_1 := T_1$) and, by the (\star)-property above, every vertex from $T_{j-1} \cap W_{j-1}$ is an endpoint of some Berge path starting in one of the sets T'_1, T''_1 . Therefore, we find a subset $T'_{j-1} \subseteq T_{j-1} \cap W_{j-1}$ with $|T'_{j-1}| \geq |T_{j-1} \cap W_{j-1}|/2 \geq \gamma m$ and $T \in \{T'_1, T''_1\}$ such that for every $v_{j-1} \in T'_{j-1}$ there is a Berge path $P = (v_1, e_1, \dots, e_{j-2}, v_{j-1})$ with $v_i \in W_i$ for all $i \in [j-2]$ such that $|e_i \cap U_{i+1}| = r-1$ for all $i \in [j-2]$ and, additionally, the vertex v_1 is from T . Again, an application of Lemma 22 allows us to extend these Berge paths such that the subset $T_j \subseteq U_j$ of all possible endpoints has cardinality at least $(1/2 + \gamma)m$. We also have $|T'_1| = \max\{1, |T_1|/2\}$ and thus: $|T'_1| = \max\{1, \gamma m/2^{j-1}\}$.

Since $\log_2(\gamma m) \leq \log_2 m - r + 1 \leq k$ the statement of the lemma follows. \square

We can now apply the above lemma iteratively, obtaining γm edge-disjoint Berge paths. The following corollary summarizes it in a ‘symmetric’ version.

Corollary 24. *For every integer $r \geq 3$, every real $\gamma \in (0, 2^{1-r})$ and $c > 1$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} on n vertices with $p \geq \frac{\ln^{cr} n}{n^{r-1}}$ and n sufficiently large. Let $H = (V, E) \subseteq \mathcal{H}$ be a*

subgraph of \mathcal{H} and $m \geq \frac{n}{\ln^c n}$. Let $U_1, U_2, \dots, U_t, \dots, U_{t+t'}$ (for some $t, t' \geq \log_2 m$) be pairwise disjoint subsets of $V(H)$ each of cardinality m . Further let A and B be two (not necessarily disjoint) sets of cardinality at least $3\gamma m$ each, disjoint from the other sets U_i . Assume moreover that the following conditions on the vertex degrees are satisfied

- (i) $\deg_H(u, U_1) \geq \left(\frac{1}{2^{r-1}} + 2\gamma\right) p\binom{m}{r-1}$ for every $u \in A$,
- (ii) $\deg_H(u, U_{t+t'}) \geq \left(\frac{1}{2^{r-1}} + 2\gamma\right) p\binom{m}{r-1}$ for every $u \in B$,
- (iii) $\deg_H(u, U_{i+1}) \geq \left(\frac{1}{2^{r-1}} + 2\gamma\right) p\binom{m}{r-1}$ for every $u \in U_i$ and every $i \in [t]$,
- (iv) $\deg_H(u, U_{t+t'-i}) \geq \left(\frac{1}{2^{r-1}} + 2\gamma\right) p\binom{m}{r-1}$ for every $u \in U_{t+t'+1-i}$ and every $i \in [t']$,

Then for any ordering of A as a_1, a_2, \dots and B as b_1, b_2, \dots , there exists a system of pairwise edge-disjoint Berge paths $P_1, \dots, P_{\gamma m}$ such that

- (a) there exists a γm -set $\{i_1, \dots, i_{\gamma m}\}$ so that the endpoints of P_j are a_{i_j} and b_{i_j} for every $j \in [\gamma m]$ and the edges lie completely within $(\cup_s U_s) \cup \{a_{i_j}, b_{i_j}\}$,
- (b) every vertex in $(\cup_s U_s)$ is an inner vertex of at most one of the paths P_j ,
- (c) every path P_j has exactly one inner vertex from each of the sets $A, B, U_1, \dots, U_{t+t'}$.

Proof. We choose $\varepsilon > 0$ so that Lemma 23 is applicable on input r, γ and c . We will apply now Lemma 23 γm times, each time finding a new Berge path. More precisely, every time we delete the inner vertices of the Berge path we just found and then continue with another round of application of Lemma 23. After γm applications, we have γm Berge paths P_i as described in (a)–(c).

Assume that we already found some paths P_1, \dots, P_s with $s < \gamma m$. Let $D_1, \dots, D_{t+t'}$ be the sets of inner vertices of these paths with $D_i \subseteq U_i$ (in the beginning we have $s = 0$ and all D_i 's are empty). To construct a new Berge path, we apply Lemma 23 twice: to some γm -set $W_0 \subset A \setminus \{a_{i_j} : j \in [s]\}$ and the sets $W_i := U_i \setminus D_i$ ($i \in [t]$) and to some γm -set $W_{t+t'+1} \subset B \setminus \{b_{i_j} : j \in [s]\}$ and the sets $W'_i := U_{t+t'-i+1} \setminus D_{t+t'-i+1}$ ($i \in [t' + 1]$). This yields vertices $v_0 \in W_0$ and $v_{t+t'+1} \in W_{t+t'+1}$ and sets $T_t \subseteq W_t$ and $T'_t \subseteq W_t$, each of size at least $(1/2 + \gamma)m$ so that for every $v_t \in T_t$ there exists a Berge path that starts in v_0 and ends in v_t and the inner vertices of which avoid already used vertices from the sets D_i , and such that for every $v_t \in T'_t$ there exists a Berge path with similar properties ending in $v_{t+t'+1}$. We may assume that $v_0 = a_j$ and $v_{t+t'+1} = b_j$ for some j , not previously used, since we can apply Lemma 22 to any γm -sets $W_0 \subset A$ and $W_{t+t'+1} \subset B$ and thus less than γm vertices of A and of B will fail to serve as end vertices of a Berge path.

Since $|T_t \cap T'_t| \geq 2\gamma m$, we finally obtain a Berge path P_j connecting a_j and b_j as required for the properties (a)–(c). Iterating this procedure (applied to W_i 's instead of U_i 's, the conditions (i)–(iv) change only negligibly) yields the desired system of γm Berge paths. \square

Next we show how iterating the above corollary will allow us to connect a_i with b_i for every i .

Lemma 25 (Connecting lemma). *For every integer $r \geq 3$, every real $\gamma \in (0, 2^{1-r})$ and $c > 7$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom r -uniform hypergraph \mathcal{H} on n vertices with $p \geq \frac{\ln^{cr} n}{n^{r-1}}$, $\Delta_2(\mathcal{H}) \leq 2 \ln n$ and n sufficiently large. Let $H = (V, E) \subseteq \mathcal{H}$ be a subgraph of \mathcal{H} , let $A = \{a_1, \dots, a_m\}$, $B = \{b_1, \dots, b_m\}$ and U be subsets of V with*

- (i) $(A \cup B) \cap U = \emptyset$,
- (ii) $m \geq \frac{2n}{\ln^c n}$,
- (iii) $|U| \geq \frac{3}{\gamma} m \log_2^2 m$,
- (iv) $\deg_H(u, U) \geq \left(\frac{1}{2^{r-1}} + 3\gamma\right) p \binom{|U|}{r-1}$ for every $u \in A \cup B \cup U$,

Then there exists a system of pairwise edge-disjoint Berge paths P_1, \dots, P_m such that

- (a) the endpoints of P_i are a_i and b_i for every $i \in [m]$ and the edges lie completely within $U \cup \{a_i, b_i\}$ (if $a_i = b_i$ we abuse notation still calling P_i a Berge path, although it is actually a Berge cycle),
- (b) every vertex in U is an inner vertex of at most one of the paths P_i ,
- (c) the length of each Berge path P_i is between $2 \log_2 m + 1$ and $4 \log_2 m + 2$, and all lengths may be chosen to be even or odd at the same time.

Proof. We choose $\varepsilon > 0$ so that Corollary 24 is applicable on input r , γ and c . Notice that we may then apply Lemma 13 on input r , γ , $c' = 2.5$ and c .

The proof strategy will proceed in rounds. In each round we will connect half of the yet not connected pairs (a_i, b_i) . It is clear, that after $\log_2 m$ rounds the process will terminate. Next we turn to the technical details. We describe the first two rounds and it will become clear how we proceed in the remaining rounds. We assume that all paths should have odd length (as the case of even length is treated similarly, by choosing an additional set $U'_{2 \log_2 m + 1}$ of size m in every round).

We come to the *first round*. We use Lemma 13 to consecutively choose pairwise disjoint subsets $U'_1, \dots, U'_{2 \log_2 m}$ of size m of the set U such that the following holds (throughout the proof, the parameter $\ell = O(\log_2^2 m)$ is the number of subsets of size m or $2m$ that we chose from U so far):

$$(\star) \quad \deg_H(u, U') \geq \left(\frac{1}{2^{r-1}} + 3\gamma - \frac{\ell}{\ln^{c'} n}\right) p \binom{|U'|}{r-1} \text{ for every } u \in \{a_1, \dots, a_m\} \cup \{b_1, \dots, b_m\} \cup U \text{ and every } U' \in \left\{U'_1, \dots, U'_{2 \log_2 m}, U \setminus \left(\bigcup_{i=1}^{2 \log_2 m} U'_i\right)\right\}.$$

It is clear, that Corollary 24 is applicable. Indeed, the degree conditions (i)–(iv) are satisfied which allows us to find γm Berge paths with properties (a)–(c) as asserted by Corollary 24 (hence, also satisfying the assertions (a)–(c) of this Lemma). We can then continue with finding each time new $2 \log_2 m$ sets with property (\star) . We do this by applying Lemma 13 (which is possible, because sampling $2 \log_2 m$ sets does not change

the parameters significantly and the lower bounds on the minimum degrees (i) and (ii) are still in place). We apply Corollary 24 $1/(2\gamma)$ times by connecting each time γm of the yet not connected pairs (a_i, b_i) through Berge paths. Thus, in the end we connect $m/2$ pairs (a_i, b_i) . Overall, we also sample in the first round (with Lemma 13) $\gamma^{-1} \log_2 m$ pairwise disjoint m -subsets of U . We delete the vertices of the sampled sets from U , but use the same notation for simplicity. This finishes the first round.

We move next to the *second round*. Let I be an $m/2$ -set of indices of pairs (a_i, b_i) which have not been yet connected. We still denote the vertices $\{a_i : i \in I\}$ by A and similarly for B . We apply Lemma 13 to choose two disjoint subsets A_1 and B_1 of U , each of cardinality $2m$ such that

$$\deg_H(a, A_1) \geq \left(1 - \frac{1}{\ln^{c'} n}\right) \left(\frac{1}{2^{r-1}} + 3\gamma - \frac{\ell}{\ln^{c'} n}\right) p \binom{|A_1|}{r-1} \text{ for all } a \in A,$$

and

$$\deg_H(b, B_1) \geq \left(1 - \frac{1}{\ln^{c'} n}\right) \left(\frac{1}{2^{r-1}} + 3\gamma - \frac{\ell}{\ln^{c'} n}\right) p \binom{|B_1|}{r-1} \text{ for all } b \in B.$$

Then Lemma 19 (applied with $m/2$ instead of m) asserts the existence of 2-matchings $(e_a, f_a)_{a \in A}$ and $(e_b, f_b)_{b \in B}$ for (A, A_1) , for (B, B_1) respectively. Let τ_A and τ_B be the injections for these 2-matchings, cf. Definition 18. Let $A' := \tau_A(A) \subseteq A_1$ and $B' := \tau(B) \subseteq B_1$. Next we order the vertices of A' according to the index set $I = \{i_1, \dots, i_{m/2}\}$ as follows: $\tau_A(e_{a_{i_1}}), \tau_A(f_{a_{i_1}}), \tau_A(e_{a_{i_2}}), \tau_A(f_{a_{i_2}}), \dots, \tau_A(e_{a_{i_{m/2}}}), \tau_A(f_{a_{i_{m/2}}})$. Exactly in the same way we order the vertices of B' . We are now back in the original situation: *exactly* as in the first round we sample at most $(2/\gamma) \log_2 m$ many pairwise disjoint m -subsets of U , and thus, we find a system of $m/2$ Berge paths $P_1^{(2)}, \dots, P_{m/2}^{(2)}$ between the sets A' and B' .

Out of $m/2$ Berge paths $P_1^{(2)}, \dots, P_{m/2}^{(2)}$ we find at least $m/4$ so that no two of them have endpoints of the form $\tau_A(e_{a_{i_j}}), \tau_A(f_{a_{i_j}})$ or $\tau_B(e_{b_{i_s}}), \tau_B(f_{b_{i_s}})$ for some j or s respectively. Recall, that we have 2-matchings for (A, A') and for (B, B') . Using these we can find at least $m/4$ Berge paths and extend these to $m/4$ Berge paths between A and B that satisfy properties (a)–(c). Observe that in this second round we sampled (with Lemma 13) disjoint subsets of U the union of which contains at most $\gamma^{-1} m \log_2 m + 4m < (\frac{2}{\gamma}) m \log_2 m$ vertices.

There remain $m/4$ vertices of A and B to be matched, and the corresponding $m/2$ vertices of A_1 and B_1 , which are connected by 2-matchings respectively. We proceed in a similar way in the next rounds by finding 2-matchings between appropriate sets, finding $m/2$ Berge paths and then identifying $m/2^i$ additional Berge paths in the i th round. After $\log_2 m$ rounds we find the desired system of m Berge paths by using at most $(2/\gamma) m \log_2^2 m$ vertices.

The paths in each round get longer by additional two inner vertices, and have length at most $4 \log_2 m + 2$ in the end. Moreover, the degree requirements of Lemma 13, Lemma 19 and Corollary 24 are reduced in each step by at most $\frac{1}{\ln^{c'} n} p \binom{|C|}{r-1}$ (where C equals U, U_i, A_1 or B_1 respectively), which guarantees their applicability throughout all rounds. This finishes the proof. \square

6 Proof of Theorem 5

6.1 Proof outline

We explain the idea of the proof for weak Hamilton cycles and not to worry about the Berge property.

Given a spanning subhypergraph $H \subseteq \mathcal{H} \sim H^{(r)}(n, p)$ as in the statement of the theorem, we partition the vertex set of H into disjoint sets Y , Z and W , where $|Y| = o(n)$, $|Z| = n/\log^{\mathcal{O}(1)} n$ and W contains almost all vertices of H . The set Z assumes the role of a reservoir. Choosing a partition with such sizes uniformly at random guarantees that with positive probability for every vertex v the edges incident to v are distributed as expected into the sets Y , Z and W (cf. Lemma 13). Also, since H is a subgraph of the random hypergraph, we have good control on the edge distribution among various subsets of vertices (cf. Lemmas 10 and 11).

Next we construct a weak path Q with $V^*(Q) \subseteq Y$ such that for every subset $M \subseteq Z$ there exists a weak path Q_M that has the same endpoints as Q and such that $V^*(Q_M) = V^*(Q) \dot{\cup} M$ (we will use Lemma 25 to construct absorbers as described by Proposition 21 and to connect them into the path Q). This property (*absorbing property*) will be crucial at a later stage of the argument.

Then we partition W randomly into $\log^{\mathcal{O}(1)} n$ sets and distribute $Y \setminus V^*(Q)$ among them such that all of these sets have the same size $o(n/\log^{\mathcal{O}(1)} n)$. Informally speaking, since $|Y|$ is significantly smaller than $|W|$ and every vertex from Y is ‘well-connected’ to W , such partition allows us to find weak paths P_1, \dots, P_m (using Lemma 16, some paths may consist of single vertices), with $m = n/\log^{\mathcal{O}(1)} n$, so that $V^*(P_1), \dots, V^*(P_m)$ form a partition of $W \dot{\cup} Y \setminus V^*(Q)$.

As a last step, we use vertices from Z to connect the paths P_1, \dots, P_m and Q into a weak cycle C (again this is possible since every vertex of H is ‘well-connected’ into Z). Since the unused vertices M of Z can be absorbed by the path Q into a weak path Q_M with $V^*(Q_M) = V^*(Q) \dot{\cup} M$, we have found a weak Hamilton cycle in H in this way. To construct the path Q and to connect the paths P_1, \dots, P_m and Q into a cycle we will repeatedly use the connecting lemma (Lemma 25) that will allow us to connect various vertices by paths of length $\mathcal{O}(\log n)$.

6.2 Rigorous details: proof of Theorem 5

We show the following result about robustness of pseudorandom hypergraphs.

Theorem 26. *For every integer $r \geq 3$ and every real $\gamma > 0$ there exists an $\varepsilon > 0$ such that the following holds for any (ε, p) -pseudorandom hypergraph \mathcal{H} on n vertices with $\Delta_2(\mathcal{H}) \leq 2 \ln n$, $p = \frac{\ln^{17r} n}{n^{r-1}}$ and n sufficiently large. Let $H \subseteq \mathcal{H}$ be a spanning subgraph with $\delta_1(H) \geq \left(\frac{1}{2^{r-1}} + \gamma\right) p \binom{n}{r-1}$. Then H contains a Hamilton Berge cycle.*

Observe first that Theorem 26 implies immediately Theorem 5 for the probability $p = \frac{\ln^{17r} n}{n^{r-1}}$, since a.a.s. the random hypergraph $H^{(r)}(n, p)$ is (ε, p) -pseudorandom by Lemmas 10 and 11 and also satisfies $\Delta_2(\mathcal{H}) \leq 2 \ln n$ a.a.s., by Proposition 9. The following

easy proposition briefly shows how the statement then extends to all $p \geq \frac{\ln^{17r} n}{n^{r-1}}$ in a straightforward way.

Proposition 27. *If Theorem 5 holds for $p = \frac{\ln^c n}{n^{r-1}}$ then it is also true for $p \geq \frac{\ln^c n}{n^{r-1}}$.*

Proof. Let $\mathcal{H} \sim H^{(r)}(n, p)$ for some $p \geq \frac{\ln^c n}{n^{r-1}}$ and let $H \subseteq \mathcal{H}$ be a subgraph of \mathcal{H} with minimum vertex degree at least $(\frac{1}{2^{r-1}} + \gamma)p \binom{n}{r-1}$. We set $q := \frac{\ln^c n}{p \cdot n^{r-1}}$ and denote by G_q the random subgraph of G where each edge is kept with probability q independently of the other edges. Clearly, $\mathcal{H}_q \sim H^{(r)}(n, qp)$ and $H_q \subseteq \mathcal{H}_q$ and, by Chernoff's inequality (Theorem 7) a.a.s. $\delta_1(H_q) \geq (\frac{1}{2^{r-1}} + \gamma/2)pq \binom{n}{r-1}$, and $\Delta_2(\mathcal{H}_q) \leq 2 \ln n$ a.a.s. as well (Proposition 9). Hence, we may apply Theorem 5 in the special case when $pq = \frac{\ln^c n}{n^{r-1}}$ and the general claim follows. \square

Proof of Theorem 26.

Setup. W.l.o.g. we assume that $\gamma < 2^{3-r}$ and we set $\gamma' = \gamma/4$, $c' := 7.4$. We thus have $p = \frac{\ln^{17r} n}{n^{r-1}}$. Let ε_1 be as asserted by Lemma 16 on input r , γ' and $c = 17$, let ε_2 be as asserted by Lemma 25 on input r , γ' and $c = 17$. We set $\varepsilon := \min\{\varepsilon_1, \varepsilon_2\}$.

We apply Lemma 13 to $V(H)$ twice to obtain three pairwise-disjoint sets W , Y and Z with the following properties:

- (i) $|Z| = \frac{n}{\log_2^4 n}$ and $\deg(v, Z) \geq (2^{1-r} + 4\gamma' - \frac{2}{\ln^{c'} n})p \binom{|Z|}{r-1}$ for all $v \in V$,
- (ii) $|Y| = \frac{9}{\gamma'} |Z| \log_2^3 n$ and $\deg(v, Y) \geq (2^{1-r} + 4\gamma' - \frac{2}{\ln^{c'} n})p \binom{|Y|}{r-1}$ for all $v \in V$,
- (iii) $|W| = V \setminus (Y \cup Z)$ and $\deg(v, W) \geq (2^{1-r} + 4\gamma' - \frac{2}{\ln^{c'} n})p \binom{|W|}{r-1}$ for all $v \in V$.

Constructing an absorbing path P_A . Our aim here is to construct absorbers for every $u \in Z$ and to put them into a single path. For this we will use vertices from Y , by applying Lemma 13 followed by Lemma 25 several times. More precisely, we construct for every $u \in Z$ a u -absorber A_u such that the inner vertices of all A_u 's are pairwise disjoint and the edges are pairwise disjoint as well. We will do it in three stages.

In the *first stage* we apply Lemma 13 to the set Y and obtain a set U_1 of cardinality $\frac{3}{\gamma'} |Z| \log_2^2 n$ such that

- (i) $\deg_H(v, U_1) \geq (2^{1-r} + 3\gamma')p \binom{|U_1|}{r-1}$ for all $v \in V$, and
- (ii) $\deg_H(v, Y \setminus U_1) \geq (2^{1-r} + 4\gamma' - \frac{3}{\ln^{c'} n})p \binom{|Y| - |U_1|}{r-1}$ for all $v \in V$.

We then apply Lemma 25 (with $A, B = Z$ and $U = U_1$) to put each vertex $u \in Z$ on its own Berge cycle C_u of *odd* length between $2 \log_2 |Z| + 1$ and $4 \log_2 |Z| + 1$, in order to obtain the cycles needed by Definition 20.

In the *second stage* we need to connect the corresponding pairs of vertices on each of the Berge cycles C_u , where $u \in Z$, as specified in the property (ii) from the definition of the asorbers, Definition 20. Observe that there are at most $2 \log_2 |Z| - 1$ such pairs for each absorber, which requires connecting in total at most $2|Z| \log_2 |Z|$ pairs. We apply Lemma 13 to the set $Y \setminus U_1$ and obtain a set U_2 of cardinality $\frac{6}{\gamma'} |Z| \log_2^3 n$ such that

- (i) $\deg_H(v, U_2) \geq (2^{1-r} + 3\gamma')p\binom{|U_2|}{r-1}$ for all $v \in V$, and
- (ii) $\deg_H(v, Y \setminus (U_1 \cup U_2)) \geq (2^{1-r} + 4\gamma' - \frac{4}{\ln^c n})p\binom{|Y| - |U_1| - |U_2|}{r-1}$ for all $v \in V$.

Again, an application of the connecting lemma, Lemma 25, yields the system of Berge paths, that completes for each $u \in Z$ an absorber A_u with the required properties. More precisely, the property (ii) of Definition 20 specifies the vertex pairs which are to be connected by Berge paths (with pairwise disjoint inner vertices), and we apply Lemma 25 to these pairs (with the set U_2 as U).

Finally, in the *third stage*, we put all our absorbers onto a Berge path. More precisely, we aim to connect the paths of all A_u 's as specified in Proposition 21 (a) into a Berge path. We consider the main endpoints u_i and u'_i of every u -absorber and we wish to connect u'_i with u_{i+1} for every $i \in [|Z| - 1]$ by edge disjoint Berge paths (whose inner vertex sets are pairwise disjoint as well), using again the connection lemma, Lemma 25. Recall, that by our choice of Y , we have $|Y \setminus (U_1 \cup U_2)| \geq \frac{4}{\gamma'}|Z| \log_2^2 n$ and $\deg_H(v, Y \setminus (U_1 \cup U_2)) \geq (2^{1-r} + 3\gamma')p\binom{|Y| - |U_1| - |U_2|}{r-1}$ for all $v \in V$. Therefore, all assumptions of Lemma 25 are met, and we obtain thus a system of Berge paths that connects the required pairs of endpoints of the absorbers A_u ($u \in Z$) into the *absorbing Berge path* P_A . Observe that $V^*(P_A) \subseteq Z \cup Y$, the endpoints of P_A are u_1 and $u'_{|Z|}$ and the path P_A has the following *absorbing property* (by Proposition 21):

- for every subset $Z' \subseteq Z$, there exists a Berge path $P_{Z'}$ with the same endpoints as P_A and $V^*(P_{Z'}) = V^*(P_A) \setminus Z'$.

We will later use some vertices from Z for further connections that come and we denote by Y' the vertices of Y which are neither inner vertices of P_A .

Partitioning W . As a next step we partition W into $\log_2^7 n$ sets W_1, \dots, W_i of the same cardinality plus a remaining set M of fewer than $n/\log_2^7 n$ vertices. We do so by consecutively applying Lemma 13, so that the sets W_i will satisfy the following properties:

- $\deg(v, W_i) \geq (2^{1-r} + 3\gamma')p\binom{|W_i|}{r-1}$ for all $v \in V$.

Next we distribute the vertices from Y' equally among the sets W_i 's (but we still use the same notation for these new sets) and put the at most $\log_2^7 n$ vertices to M . Observe that then still $\deg(v, W_i) \geq (2^{1-r} + 2\gamma')p\binom{|W_i|}{r-1}$ for all $v \in V$ holds.

Covering W with Berge paths. Now we apply Lemma 16 to obtain (W_i, W_{i+1}) -matchings in H for $i = 1, \dots, \log_2^7 n$. Notice that this gives rise to a system of $(1 - o(1))n/\log_2^7 n$ edge-disjoint Berge paths, each of length $\log_2^7 n$, so that the inner vertices of these paths are pairwise disjoint and form a partition of $W \setminus M$. To obtain a particular path, one starts with some vertex $w_1 \in W_1$, then follows the matching edge $e_1 \ni w_1$, then considers the second endpoint $w_2 \in e_1$, then follows the matching edge $e_2 \ni w_2$ from the (W_2, W_3) -matching in H and so on.

Additionally, we also view every single vertex $w \in M$ as a Berge path. Thus, in total we cover $W \cup Y'$ by t Berge paths P_1, \dots, P_t , where t is less than $2(n/\log_2^7 n) + \log_2^7 n$.

Obtaining a weak Hamilton cycle in H . Recall, that $|Z| = \frac{n}{\log_2^4 n}$ and $\deg(v, Z) \geq (2^{1-r} + 3\gamma')p \binom{|Z|}{r-1}$ for all $v \in V$ holds. Thus, the assumptions of the connecting lemma, Lemma 25, are met. Therefore we are able to connect the Berge paths P_1, \dots, P_t and the absorbing Berge path P_A into some ‘almost’ cycle C' by using the vertices of some subset Z' from Z . By the absorbing property, we may delete the vertices from Z' and restructure the path P_A (which is a subgraph of C') into P' so that the so obtained cycle C is indeed a weak Hamilton cycle.

Why is C already a Hamilton Berge cycle in H . The cycle which we constructed in the way above is indeed Berge. The reason is that we build our cycle by constructing Berge paths between partite sets (cf. Lemmas 16, 17 and 25) and in doing so we use only edges e between some two sets U_i and U_{i+1} which lie within $U_i \cup U_{i+1}$ such that $|e \cap U_i| \in \{1, r-1\}$. In this way we guarantee that at any stage of our construction we are using genuinely new edges. Thus, the constructed cycle is indeed Berge. \square

7 Berge Hamiltonicity in dense hypergraphs

In this section we prove Proposition 3. The asymptotic tightness of the bound on the minimum vertex degree was considered already in the introduction. We remark that optimal bound was proven in a long paper in [8], whereas the proof below is elementary and very short.

Proof of Proposition 3. Let $r \geq 3$ and let $H = (V, E)$ be an r -uniform hypergraph on $n > 2r - 2$ vertices with $\delta_1(H) \geq \binom{\lceil n/2 \rceil - 1}{r-1} + n - 1$. We observe first that due to the large minimum vertex degree, H is connected. Let $P = (v_1, e_1, v_2, \dots, e_{k-1}, v_k)$ be a longest Berge path in H .

For every $v \in V$ we define $E'(v) = \{e \in E \setminus \{e_1, \dots, e_{k-1}\} : v \in e\}$. The condition on the minimum vertex degree implies that we have $|E'(v_1)|, |E'(v_k)| \geq \binom{\lceil n/2 \rceil - 1}{r-1}$. Since P is a longest Berge path, it holds for every $e \in E'(v_1)$ that $e \subseteq V^*(P)$. The same is true for v_k .

We claim that there exist distinct hyperedges $e \in E'(v_1)$ and $e' \in E'(v_k)$ as well as an index $i \in [k-1]$ such that $v_{i+1} \in e \cap V^*(P)$ and $v_i \in e' \cap V^*(P)$. Assume for a contradiction that this is not true. Then, by the pigeonhole principle, there exists a subset $S \subseteq [k-1]$ such that $|S| \leq \lfloor (k-1)/2 \rfloor$ with $f \subseteq \{v_{i+1} : i \in S\} \cup \{v_1\}$ for every $f \in E'(v_1)$ or with $f' \subseteq \{v_i : i \in S\} \cup \{v_k\}$ for every $f' \in E'(v_k)$. Suppose that $f \subseteq \{v_{i+1} : i \in S\} \cup \{v_1\}$ for every $f \in E'(v_1)$ holds. Then $\delta_1(v_1) \leq \binom{|S|-1}{r-1} + k - 1 < \binom{\lceil n/2 \rceil - 1}{r-1} + n - 1$, which is a contradiction.

Hence, there exist $e \in E'(v_1)$ and $e' \in E'(v_k)$ with the claimed property. Let

$$C = (v_1, e, v_{i+1}, e_{i+1}, v_{i+2}, \dots, v_k, e', v_i, e_{i-1}, \dots, e_1)$$

be the Berge cycle that can be constructed from P using e and e' . If $k = n$, then C is a Hamilton Berge cycle and we are done. Otherwise we get a contradiction, similar as in Dirac’s original proof, by breaking up the cycle C and extending it by a new edge (since H is connected), thus obtaining a longer Berge path. \square

Acknowledgement

We would like to thank the referee for careful reading, detailed comments and useful suggestions.

References

- [1] P. Allen, J. Böttcher, J. Ehrenmüller, and A. Taraz. The bandwidth theorem in sparse graphs. *Advances in Combinatorics*, 2020:60, 2020.
- [2] D. Bal, R. Berkowitz, P. Devlin, and M. Schacht. Hamiltonian Berge cycles in random hypergraphs. [arXiv:1809.03596](https://arxiv.org/abs/1809.03596), 2018.
- [3] J. Balogh, B. Csaba, and W. Samotij. Local resilience of almost spanning trees in random graphs. *Random Struct. Algorithms*, 38(1-2):121–139, 2011.
- [4] J. Balogh, C. Lee, and W. Samotij. Corrádi and Hajnal’s theorem for sparse random graphs. *Combin. Probab. Comput.*, 21(1-2):23–55, 2012.
- [5] J. C. Bermond, A. Germa, M. C. Heydemann, and D. Sotteau. Hypergraphes hamiltoniens. Problèmes combinatoires et théorie des graphes, Orsay 1976, Colloq. int. CNRS No. 260, 39-43 (1978)., 1978.
- [6] J. Böttcher, Y. Kohayakawa, and A. Taraz. Almost spanning subgraphs of random graphs after adversarial edge removal. *Combin. Probab. Comput.*, 22(05):639–683, 2013.
- [7] D. Clemens, J. Ehrenmüller, and Y. Person. A Dirac-type theorem for Hamilton Berge cycles in random hypergraphs. In *Discrete mathematics days. Extended abstracts of the 10th “Jornadas de matemática discreta y algorítmica” (JMDA), Barcelona, Spain, July 6–8, 2016*, pages 181–186. Amsterdam: Elsevier, 2016.
- [8] M. Coulson and G. Perarnau. A rainbow Dirac’s theorem. [arXiv:1809.06392](https://arxiv.org/abs/1809.06392), 2018.
- [9] G. A. Dirac. Some theorems on abstract graphs. *Proceedings of the London Mathematical Society*, 3(1):69–81, 1952.
- [10] A. Ferber and L. Hirschfeld. Co-degrees resilience for perfect matchings in random hypergraphs. [arXiv:1908.01435](https://arxiv.org/abs/1908.01435), 2019.
- [11] A. Ferber, R. Nenadov, A. Noever, U. Peter, and N. Škorić. Robust Hamiltonicity of random directed graphs. *J. Comb. Theory, Ser. B*, 126:1–23, 2017.
- [12] D. Hefetz, A. Steger, and B. Sudakov. Random directed graphs are robustly Hamiltonian. *Random Struct. Algorithms*, 49(2):345–362, 2016.
- [13] H. Huang, C. Lee, and B. Sudakov. Bandwidth theorem for random graphs. *J. Comb. Theory Ser. B*, 102(1):14–37, 2012.
- [14] S. Janson, T. Łuczak, and A. Ruciński. *Random graphs*, volume 45. John Wiley & Sons, 2011.
- [15] M. Krivelevich, C. Lee, and B. Sudakov. Resilient pancyclicity of random and pseudorandom graphs. *SIAM J. Discrete Math.*, 24(1):1–16, 2010.

- [16] D. Kühn and D. Osthus. Hamilton cycles in graphs and hypergraphs: an extremal perspective. In *Proceedings of the International Congress of Mathematicians, Seoul, Korea*, volume 4, pages 381–406, 2014.
- [17] C. Lee and B. Sudakov. Dirac’s theorem for random graphs. *Random Struct. Algorithms*, 41(3):293–305, 2012.
- [18] R. Montgomery. Hamiltonicity in random directed graphs is born resilient. [arXiv:1901.09605](https://arxiv.org/abs/1901.09605), 2019.
- [19] R. Montgomery. Hamiltonicity in random graphs is born resilient. *J. Comb. Theory, Ser. B*, 139:316–341, 2019.
- [20] R. Nenadov, A. Steger, and M. Trujić. Resilience of perfect matchings and Hamiltonicity in random graph processes. *Random Struct. Algorithms*, 54(4):797–819, 2019.
- [21] D. Poole. On weak Hamiltonicity of a random hypergraph. [arXiv:1410.7446](https://arxiv.org/abs/1410.7446), 2014.
- [22] V. Rödl and A. Ruciński. Dirac-type questions for hypergraphs — a survey (or more problems for Endre to solve). In *An Irregular Mind*, pages 561–590. Springer, 2010.
- [23] V. Rödl, A. Ruciński, and E. Szemerédi. A Dirac-type theorem for 3-uniform hypergraphs. *Combin. Probab. Comput.*, 15(1-2):229–251, 2006.
- [24] N. Škorić, A. Steger, and M. Trujić. Local resilience of an almost spanning k -cycle in random graphs. *Random Structures & Algorithms*, 53(4):728–751, 2018.
- [25] B. Sudakov and V. H. Vu. Local resilience of graphs. *Random Struct. Algorithms*, 33(4):409–433, 2008.
- [26] Y. Zhao. Recent advances on Dirac-type problems for hypergraphs. In *Recent trends in combinatorics*, pages 145–165. Cham: Springer, 2016.